

THE CORE-STABLE SETS-THE BARGAINING THEORY FROM A FUNCTIONAL  
AND MULTI CRITERION VIEWPOINT

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Abstract:

In this paper, it is introduced the concept of  $f$ -imputation from which the core is defined and a theorem of analogous characterization to that given in Owen(1982) is proved.

Also, the bargaining theory from the viewpoint analogous to that developed in Davis and Maschler(1963) and Peleg(1963) is exposed.

### 1. Introduction

In his excellent book [2], G.Owen provides a characterization of the core of a game as a subset of  $\mathbb{R}^n$ . There, he defines the usual notions of imputations and domination for cooperative  $n$ -person games. Davis-Maschler and Peleg in [1] and [3], introduce the notion of stable coalitions, bargaining sets and prove existence theorems for the bargaining set  $M_1^{(1)}$  in euclidean spaces.

In this paper we introduce the concept of  $f$ -imputation which generalizes the classical notion of imputation. We also extend the concept of core. In particular, we characterize the latter as a subset of a topological space. Besides, in the same framework, we study the bargaining set and prove an existence theorem only assuming the continuity of the function  $f$ .

$X$  will indicate a compact connected subset of a topological space.

$N$  will indicate a finite set of index,  $\text{card}(N)=n$ .

$v$  will indicate a defined function on the subsets of  $N$  to nonnegative real values such that:

$$v(\emptyset)=0 \tag{1-1}$$

$$v(S \cup T) \geq v(S) + v(T) \quad , \quad S \cap T = \emptyset \tag{1-2}$$

For all  $i \in N$  let  $f_i : X \rightarrow [0, \infty)$  and we indicate  $f : X \rightarrow [0, \infty)^N$  to the application defined by

$$f(x) = \{f_i(x)\}_{i \in N}$$

Definition 1-1: An element  $x \in X$  is an  $f$ -imputation for a game  $v$ , if :

$$\left\{ \begin{array}{l} \text{i) } \sum_{i \in N} f_i(x) = v(N) \\ \text{ii) } f_i(x) \geq v(\{i\}) \quad \text{for all } i \in N \end{array} \right. \quad (1-3)$$

Definition 1-2: Let  $x$  and  $y$  be two  $f$ -imputations,  $S \subset N$ , then we say that  $x$  dominates  $y$  through  $S$  and we denote this by  $x \succ_S y$ , if

$$\left\{ \begin{array}{l} \text{i) } f_i(x) > f_i(y) \quad \text{for all } i \in S \\ \text{ii) } \sum_{i \in S} f_i(x) \leq v(S) \end{array} \right. \quad (1-4)$$

## 2 - The Core

Definition 2-1: The set of all undominated  $f$ -imputations for a game  $v$ , will be called core and we will denote it by  $C(v)$

Theorem 2.1: Let  $f : X \rightarrow [0, \infty)^N$  be surjective, then the core for game  $v$  is the set of all  $x \in X$  that satisfy:

$$\left\{ \begin{array}{l} \text{i) } \sum_{i \in S} f_i(x) \geq v(S) \quad \text{for all } S \subset N \\ \text{ii) } \sum_{i \in N} f_i(x) = v(N) \end{array} \right. \quad (2-1)$$

Proof:

Let  $x$  be (2-1) i) and ii)

If  $S = \{i\}$  the condition i) means that  $f_i(x) \geq v(\{i\})$  that together with the condition ii) means that  $x$  is an  $f$ -imputation.

$x$  is undominated, in fact, let us suppose that there exists  $y \in X$  and  $S \subset N$  such that  $f_i(y) > f_i(x)$  for all  $i \in S$ , but this together with (2-1) i) means

$$\sum_{i \in S} f_i(y) > v(S)$$

and this contradicts (1-4) i). Hence  $x \in C(v)$ .

Conversely, suppose that  $y$  does not satisfy (2-1) i) or ii).

If ii) fails,  $y$  is not an  $f$ -imputation and hence  $y \notin C(v)$ .

If  $y$  is such that it does not verify i) then there exists  $S \subset N$  such that

$$\sum_{i \in S} f_i(y) < v(S); \text{ this is } \sum_{i \in S} f_i(y) = v(S) - \epsilon \text{ with } \epsilon > 0.$$

$$\text{Let } \alpha = v(N) - v(S) - \sum_{i \in S} v(\{i\}) \text{ and } \Delta = \text{card}(S) - \alpha \geq 0.$$

Let  $t = \{t_i\}_{i \in N} \in [0, \infty)^n$  where

$$t_i = \begin{cases} f_i(y) + \frac{\epsilon}{\Delta} & \text{if } i \in S \\ v(\{i\}) + \frac{\alpha}{n - \Delta} & \text{if } i \notin S \end{cases}$$

then by the surjectivity of  $f$ , there exists  $z \in X$  such that  $f(z) = t$ , then :

$$f_i(z) = \begin{cases} f_i(y) + \frac{\epsilon}{\Delta} & \text{if } i \in S \\ v(\{i\}) + \frac{\alpha}{n - \Delta} & \text{if } i \notin S \end{cases}$$

Clearly  $z$  is an  $f$ -imputation and  $z \succ_S y$ , then  $y \notin C(v)$ . ■

### 3-The Bargaining Theory

Hence forth, let us suppose that  $v: \mathcal{P}(N) \rightarrow [0, 1]$  is such that

$$\begin{cases} \text{i) } v(\{i\}) = 0 \\ \text{ii) } v(N) = 1 \end{cases} \tag{3-1}$$

moreover properties (1-1) and (1-2).

For each  $i \in N$ ,  $f_i: X \rightarrow [0, 1]$  is continuous and  $f: X \rightarrow [0, 1]^n$  is surjective.

**Definition 3-1:** By an  $f$ -coalition structure (f.c.s.) for  $N = \{1, 2, \dots, n\}$  we shall mean a partition

$\mathcal{T} = \{T_1, T_2, \dots, T_m\}$  of  $N$

Definition 3-2 : An  $f$ -payoff configuration (f.p.c.) for a game  $v$  is :  
 $(x; \mathcal{T}) = (f_1(x), \dots, f_n(x); T_1, \dots, T_m)$ , where  $\mathcal{T}$  is an  $f$ -coalition structure (f.c.s.) and  $x \in X$  is such that

$$\sum_{i \in T_k} f_i(x) = v(T_k) \quad \text{for } k = 1, 2, \dots, m$$

Definition 3-3: Given a  $f$ -payoff configuration as in definition 3-2, we say that it is *individually rational* (i.r.f.p.c.) for a game  $v$  if it verifies that

$$f_i(x) \geq v(\{i\}) = 0 \quad \text{for all } i \in N$$

$y$  is *coalitionally rational* (c.r.f.p.c.) for a game  $v$  if verifies that

$$\sum_{i \in S} f_i(x) \geq v(S) \quad \text{for } S \subset T_k \in \mathcal{T}$$

Definition 3-4 : Let  $(x; \mathcal{T})$  be a c.r.f.p.c. for a game  $v$  and let  $\mu$  and  $\lambda$  ( $\mu \neq \lambda$ ) be belonging to an  $f$ -coalition  $T_j$  of  $\mathcal{T}$ .

An  $f$ -objection of  $\lambda$  against  $\mu$  in  $(x; \mathcal{T})$  is a vector  $f^C(y) = (f_k(y))_{k \in C}$  where  $C$  is an  $f$ -coalition containing  $\lambda$  but not  $\mu$ , and where its coordinates satisfy :

$$\begin{aligned} & \text{and} \quad f_\lambda(y) > f_\lambda(x) \\ & \text{and} \quad f_k(y) \geq f_k(x) \quad (k \neq \lambda; k \in C) \\ & \sum_{k \in C} f_k(y) = v(C) \end{aligned}$$

Definition 3-5: As in definition 3-4, an  $f$ -counter objection to this  $f$ -objection is a vector  $f^D(z) = (f_k(z))_{k \in D}$ , where  $D$  is an  $f$ -coalition containing  $\mu$  but not  $\lambda$  and whose coordinates satisfy

$$\begin{aligned} & \text{and} \quad f_k(z) \geq f_k(x) \quad \text{for each } k \in D \\ & \text{and} \quad f_k(z) \geq f_k(y) \quad \text{for each } k \in D \cap C \\ & \sum_{k \in D} f_k(z) = v(D) \end{aligned}$$

Definition 3-6 : We say that  $i$  is stronger than  $k$  (or equivalently, that  $k$  is weaker than  $i$ ) in  $(x; \mathcal{T})$  if  $i$  has an  $f$ -objection against  $k$  which cannot be  $f$ -countered.

We denote this by  $i \gg k$ . We say that  $i$  and  $k$  are equal if neither  $i \gg k$  nor  $k \gg i$ . We denote this by  $i \sim k$ .

Remark : By definition  $i \sim k$  in  $(x; \mathcal{T})$  if  $i$  and  $k$  belong to different  $f$ -coalitions.

Definition 3-7: An  $f$ -coalition  $T_j$  in  $\mathcal{T}$  is called  $f$ -stable in  $(x; \mathcal{T})$  if each two of its members are equal.

Definition 3-8 : The set of all  $f$ -stable individually rational  $f$ -payoff configurations is called the  $f$ -bargaining set and we denote it by  $M_1^{(1)}(f)$ . Given an  $f$ -coalition structure  $\mathcal{T}$ , we denote  $X(\mathcal{T})$  the set of  $x \in X$  such that  $(x; \mathcal{T})$  is an i.r.f.p.c.

Lemma 3-1 : Let  $c_1(x), c_2(x), \dots, c_n(x)$  be continuous functions defined for  $x \in X(\mathcal{T})$  to nonnegative real values.

If, for each  $x \in X(\mathcal{T})$  and for each  $T_j \in \mathcal{T}$  there exists  $i \in T_j$  such that  $c_i(x) \geq f_i(x)$  then, there exists  $\xi \in X(\mathcal{T})$  such that  $c_i(\xi) \geq f_i(\xi)$  for each  $i \in N$ .

Proof:

For  $x \in X(\mathcal{T})$  and  $i \in N$  we denote, using the surjectivity of  $f$ ,

$$f_i(z) = \begin{cases} f_i(x) - c_i(x) & \text{if } f_i(x) \geq c_i(x) \\ 0 & \text{if } f_i(x) < c_i(x) \end{cases} \quad (3-2)$$

and if  $i \in T_j$

$$f_i(y) = f_i(x) - f_i(z) + \frac{1}{\tau_j} \sum_{k \in T_j} f_k(z) \quad (3-3)$$

where  $\tau_j = \text{card}(T_j)$

It is clear that  $f(y)$  is a continuous function of  $f(x)$ . Moreover, it can be seen that  $f_i(y) \geq 0$  and  $\sum_{i \in T_j} f_i(y) = v(T_j)$  and as  $0 = v(\{i\}) \leq f_i(y)$  then  $y \in X(\mathcal{T})$ .

Let us suppose now  $f_i(x) > c_i(x)$ . This means that  $f_i(z) > 0$ . Moreover, there exists  $k \in T_j$  such that  $f_k(x) \leq c_k(x)$ , then by (3-2),  $f_k(z) = 0$ . Hence

$$f_k(y) \geq f_k(x) + \frac{f_1(z)}{\tau_j} > f_k(x)$$

then  $f(x)$  is not a fixed point by the application of  $[0,1]^n$  in  $[0,1]^n$  that to  $f(x)$  it assigns  $f(y)$  defined in (3-3). Then, by Brouwer's fixed point theorem, there exists  $\xi \in X(\mathcal{J})$  such that

$$f_1(\xi) = f_1(\xi) - f_1(z) + \frac{1}{\tau_j} \sum_{k \in I_j} f_k(z)$$

and clearly, this means by (3-2) that

$$f_1(\xi) \leq c_1(\xi) \quad \text{for all } i \in N.$$

**Definition 3-9** : Let  $(x; \mathcal{J})$  be an i.r.f.p.c., and let  $C$  be an  $f$ -coalition. Then the  $f$ -excess of  $C$  is

$$e(C) = v(C) - \sum_{i \in C} f_i(x)$$

**Lemma 3-2** : If in  $(x; \mathcal{J})$ ,  $\lambda$  has an  $f$ -objection  $f^C(y)$  against  $\mu$  and this  $f$ -objection cannot be  $f$ -countered, then each  $f$ -coalition  $D$ , for  $\mu \in D$ , and  $e(D) \geq e(C)$ , must contain  $\lambda$ .

**proof:**

Let us suppose that  $e(D) \geq e(C)$  and  $\lambda \notin D$  we shall see that there exists  $z \in X$  such that  $f^D(z)$  is an  $f$ -counter objection of  $\mu$  against  $\lambda$ .

Let  $z \in X$ , such that

$$f_k(z) = \begin{cases} f_k(y) & \text{if } k \in C \cap D \\ f_k(x) + \varepsilon_k & \text{if } k \in D - C \end{cases} \quad (3-4)$$

We compute  $\varepsilon_k \geq 0$

In fact, by hypothesis:

$$v(D) - v(C) + \sum_{k \in D-C} f_k(x) - \sum_{k \in D-C} f_k(x) \geq 0 \quad (3-5)$$

and

$$v(D) = v(C) - \sum_{k \in C-D} f_k(x) + \sum_{k \in D-C} f_k(x) + \sum_{k \in D-C} \varepsilon_k$$

Then, by (3-5)

$$\sum_{k \in D-C} \varepsilon_k = v(D) - v(C) + \sum_{k \in C-D} f_k(x) - \sum_{k \in D-C} f_k(x) \geq 0$$

Selecting

$$\epsilon_k = \frac{v(D) - v(C) + \sum_{C \rightarrow D} f_k(x) - \sum_{D \rightarrow C} f_k(x)}{\text{card}(D-C)} \geq 0$$

there results that  $f^D(z)$  is an  $f$ -counter objection.

■

**Lemma 3-3:** Let  $(x; \mathcal{T})$  be an i.r.f.p.c. Then, the relation  $\gg$  is acyclic.

proof:

It is clear that if  $i$  and  $k$  are in different  $f$ -coalitions, then  $i \sim k$ .

Let us suppose that an  $f$ -coalition  $T_1 \in \mathcal{T}$  is such that  $T_1 = \{1, 2, \dots, t\}$  and that  $1 \gg 2 \gg 3 \gg \dots \gg t$ .

Then each  $i \in T_1$  has an  $f$ -objection through the  $f$ -coalition  $C$  against  $i+1 \pmod{t}$ , which cannot be  $f$ -counter objected.

Let  $C_{i_0}$  be  $f$ -coalition (among  $C_1, \dots, C_t$ ) which has maximal  $f$ -counter objected.

We claim that  $i_0$  can  $f$ -counter object against  $i_0 - 1 \pmod{t}$  through the  $f$ -coalition  $C_{i_0}$ . Clearly  $i_0 - 1 \pmod{t}$  has only the amount  $e(C_{i_0 - 1})$  at his disposal to from the  $f$ -objecting coalition; having  $i_0$  the amount  $e(C_{i_0}) \geq e(C_{i_0 - 1})$  at his disposal, can always  $f$ -counter object unless  $i_0 - 1 \pmod{t} \in C_{i_0}$ .

Repeating this argument, we must have  $i_0 - 2 \pmod{t} \in C_{i_0}$ , etc., and eventually  $i_0 + 1 \pmod{t} \in C_{i_0}$ . But this is obviously impossible.

■

**Theorem 3-1 :** Given  $v$  as in (3-1), and  $\mathcal{T}$  any  $f$ -structure coalition. Then there exists at least  $x \in X$  such that  $(x; \mathcal{T}) \in \mathcal{M}_1^{(1)}(f)$ .

proof:

Let  $(x; \mathcal{T})$  be an i.r.f.p.c.

We denote by  $(y^J, x^{N-T_j}; \mathcal{T})$  the i.r.f.p.c. which is obtained by keeping  $f_1(x)$  fixed for  $i \in N - T_j$  and replacing  $f_k(x)$  by  $f_k(y)$  for  $k \in T_j$  where  $f_k(y) \geq 0$  and  $\sum_{k \in T_j} f_k(y) = v(T_j)$ .

Let  $E_j^1(x)$  be the set of points  $y^j$  such that in the i.r.f.p.c.  $(y^j, x^{N-T_j}, \mathcal{T})$ ,  $i$  ( $i \in T_j$ ) is not weaker than any other  $j \in N$ . The set  $E_j^1(x)$  is closed and contains the set of  $y$  from the face  $f_i(y)=0$  of simplex  $\Delta_j$  (since, if  $f_i(y)=0$ ,  $i$  can  $f$ -counter object with an  $f$ -coalition of only one element).

We define the function

$$c_i(x) = f_i(x) + \max_{y^j \in E_j^1(x)} \min_{k \in T_j} (f_k(x) - f_k(y)) \quad (3-6)$$

where  $T_j$  is the  $f$ -coalition in  $\mathcal{T}$  that contains  $i$ . It can be easily seen that  $c_i(x)$  is continuous as function of  $x$ ; since  $E_j^1(x)$  is upper and lower semi-continuous.

$E_j^1(x)$  is upper semi-continuous since given  $x_n \rightarrow x$ ;  $y_n \rightarrow y$  with  $y_n^j \in E_j^1(x_n)$ .

For each  $y_n^j \in E_j^1(x_n)$  in each i.r.f.p.c.  $(y_n^j, x_n^{N-T_j}; \mathcal{T})$   $i$  is not weaker than any other  $j \in N$ , i.e.,  $i$  has an  $f$ -objection  $f^C(z)$  against each  $j \in N$  which cannot be  $f$ -counter objected. Then

$$\begin{aligned} f_i(z) &> f_i(y_n) \\ f_k(z) &\geq f_k(y_n) && \text{for } k \in C \cap T_j \\ \sum_{k \in C} f_k(z) &= v(C) \end{aligned}$$

and for all  $f^D(t)$  where  $D$  is any  $f$ -coalition such that  $i \notin D$

$$\begin{aligned} & f_k(t) < f_k(y_n) && \text{for some } k \in D \\ \text{or} & & & \\ & f_k(t) < f_k(z) && \text{for some } k \in D \cap C \\ \text{or} & & & \\ & \sum_{k \in C} f_k(t) \neq v(D) \end{aligned}$$

Then considering the continuity of  $f_k$ , there results  $y^j \in E_j^1(x)$  and  $E_j^1$  is upper semi-continuous.

$E_j^1(x)$  is lower semi-continuous. In fact, let us suppose  $x_n \rightarrow x$ , and for all sequence  $y_n \rightarrow y$  there exists  $\eta_0$  such that  $y_{\eta_0}^j \notin E_j^1(x_{\eta_0})$ . We shall prove that



$$y^j \notin E_j^1(x).$$

By the assumption, there exists  $\mu \in T_j$  and  $f^j(z)$  such that

$$f_\mu(z) > f_\mu(y_{\eta_0})$$

and

$$f_k(z) \geq f_k(y_{\eta_0}) \quad \text{for } k \in T_j$$

then, by the continuity of  $f_k$ , there results  $y_{\eta_0}^j \notin E_j^1(x_{\eta_0})$  and  $E_j^1(x)$  is lower semi-continuous. Moreover, it can be seen that  $c_1(x)$  is nonnegative. Then, by Lemma 3-3, for any  $x \in X(\mathcal{T})$  and any  $T_j \in \mathcal{T}$ , there exists  $i \in T_j$  such that  $i$  is not weaker than any  $k \in T_j$ ; then

$$x^j \in E_j^1(x) \quad \text{and} \quad c_1(x) \geq f_1(x)$$

Then, by Lemma 3-1 there exists  $\xi$  such that  $c_1(\xi) \geq f_1(\xi)$  for all  $i \in N$ . Moreover, it is clear that

$$v(T_j) = \sum_{k \in T_j} f_k(\xi) = \sum_{k \in T_j} f_k(y), \quad \text{and } c_1(\xi) \leq f_1(\xi) \text{ for all } i, \text{ since, if}$$

there exists  $i_0 \in N$  such that  $c_{i_0}(\xi) > f_{i_0}(\xi)$ , then

$$\text{Max}_{y^j \in E_j^1(\xi)} \text{Min}_{k \in T_j} (f_k(\xi) - f_k(y)) > 0$$

and there exists  $y^j \in E_j^1(\xi)$  such that for all  $k \in T_j$ ,  $f_k(\xi) > f_k(y)$ , then

$$\sum_{k \in T_j} f_k(\xi) > \sum_{k \in T_j} f_k(y)$$

which contradicts (3-7). Hence, there results  $c_1(\xi) = f_1(\xi)$  for all  $i$ .

But this means that there exists  $y \in E_j^1(\xi)$  for all  $i$ , such that  $f_k(y) = f_k(\xi)$  and therefore  $\xi^j \in E_j^1(\xi)$  for each  $i$  and each  $j$ .

Then, in  $(\xi; \mathcal{T})$  no member is stronger than another. This means that  $(\xi; \mathcal{T}) \in \mathcal{M}_1^{(1)}$



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